

# BAPC 2019

## Solutions presentation

# A: Appeal to the Audience

Problem Author: Ragnar Groot Koerkamp

- Place teams in a tournament bracket such that the amount of skill on display to the audience is maximized.
- The best team must play the most games, 2nd the 2nd-most, etc.
- For each node, find the longest path from each child to a leaf recursively.
- Extend the longest of these path to the current node.
- Sort both the list of path lengths and the skill levels.
- Match these lists one to one for a maximal solution.

Statistics: 15 submissions, 3 accepted, 8 unknown

## B: Breaking Branches

Problem Author: Timon Knigge



Given a branch of length  $n$ , determine who will win the game if they break it into pieces repeatedly.

- A branch of length  $n$  can be cut in exactly  $n - 1$  places.
- After any cut, the number of remaining possible cuts decreases by exactly one.
- Alice wins when  $n$  is even. She can break it at any position.

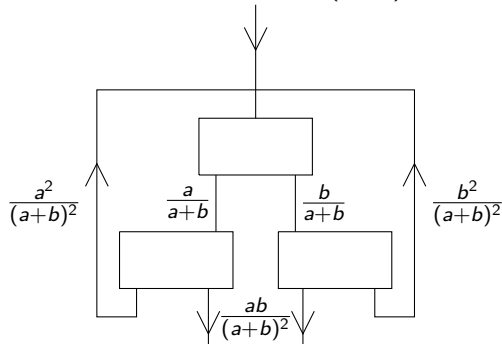
Statistics: 57 submissions, 56 accepted, 0 unknown

# C: Conveyor Belts

Problem Author: Daan van Gent

Given splitters that split their input producing an output ratio between their outputs of  $(a : b)$ , can you build a network of splitters that produces an output ratio  $(c : d)$ ?

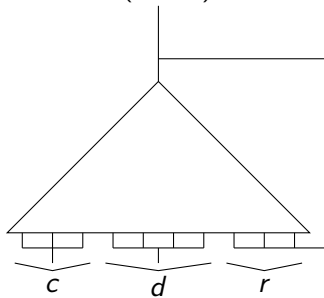
Step 1: Produce a ratio of  $(1 : 1)$ .



## C: Conveyor Belts

Problem Author: Daan van Gent

Step 2: Create a binary tree of depth  $n$  of these (1 : 1) splitters, with  $n$  such that  $c + d \leq 2^n$ . Then connect  $c$  of the leafs to output 1,  $d$  leafs to output 2, and  $r = 2^n - (c + d)$  back to the root.



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## C: Conveyor Belts

Problem Author: Daan van Gent

Step 3: Remove all (1 : 1) splitters whose outputs are the same. To not run foul of timelimits, this needs to be done during generation of the tree.

Statistics: 1 submissions, 1 accepted, 0 unknown

# D: Deck Randomisation

Problem Author: Ragnar Groot Koerkamp



Given two permutations  $A$  and  $B$ , how often do we need to repeat them one after the other to get back to where we started. Two possibilities:

- Option 1:  $(AB)^n = 1$ , minimum of  $2n$
- Option 2:  $(AB)^n A = 1$ , minimum of  $2n + 1$ .

Both take similar strategy.

# D: Deck Randomisation

Problem Author: Ragnar Groot Koerkamp



First, calculate the permutation  $AB$ , and split it into cycles. Example (notation from problem):

$$A = 5\ 1\ 6\ 3\ 2\ 4$$

$$B = 4\ 6\ 5\ 1\ 3\ 2$$

$$AB = 3\ 4\ 2\ 5\ 6\ 1$$

Then  $AB$  has cycle  $(1\ 3\ 2\ 4\ 6\ 5)$



For option 1:

- A cycle of length  $k_i$  implies  $k_i | n$ .
- Hence  $n = \text{lcm}(k_i)$ .
- Example: 1 cycle of length 6, so  $n = 6$ , giving 12 shuffles.

For option 2:

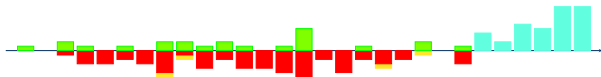
- Per cycle of length  $k_i$ , check if exists  $z$  such that  $(AB)^z = A$ . Then  $n = -z(\text{mod } k_i)$ .
  - Example:  $(AB)^4 = 516324$ , so  $n = -4(\text{mod } 6)$ .
- Option 2 is not applicable if such  $z$  does not exist.
- Then reconstruct  $n$  using the Chinese Remainder Theorem.
  - Example:  $n = 2$ , giving 5 shuffles.

Minimum of the two is answer.

Statistics: 26 submissions, 1 accepted, 11 unknown

# E: Efficient Exchange

Problem Author: Raymond van Bommel



Compute the minimal number of coin exchanges needed to pay  $n$ .

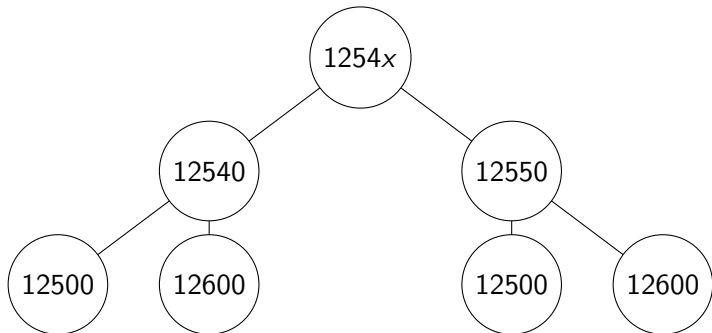
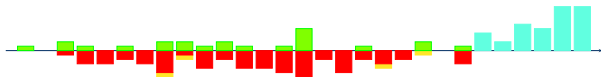
- Note that you never need to exchange more than 9 coins of value  $10^k$ , because for 10 exchanges we just use a single  $10^{k+1}$  coin.
- Consider  $1254x$ . If we can do 12540 in  $a$  exchanges and 12550 in  $b$ , then we can do

$$\min\{a + x, b + 10 - x\}.$$

- Solving recursively is costs  $2^{1000}$  calls.

# E: Efficient Exchange

Problem Author: Raymond van Bommel

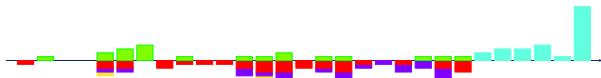


There are only two distinct nodes per layer. So only 1000 calls needed.

Statistics: 124 submissions, 23 accepted, 37 unknown

# F: Find my Family

Problem Author: Bjarki Ágúst Guðmundsson

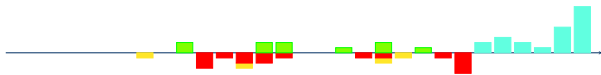


- Given a family picture, find if there are three people of relative height 2, 1, and 3, in this order.
- For each position, find the largest value on the right of it.
- Go from left to right and keep a set of all values seen so far.
- For each value, find the smallest element of the set that is larger.
- A 213 ordering exists if this smallest larger element on the left is smaller than the largest element on the right.

Statistics: 100 submissions, 20 accepted, 27 unknown

# G: Gluttonous Goop

Problem Author: Mees de Vries



- Your fungus is growing. How many squares does it occupy after  $k$  steps?
- Obvious solution: flood fill. But  $k$  can be  $10^6$ : too slow.
- We need to be smarter. Let's look at an example.

Number of children	Frequency
0	0
1	1
2	2
3	3
4	10
5	8
6	5
7	3
8	2
9	1
10	1

```

#
##
#.#
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#####.
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#.#
#.#
#####.
#.#
##
###
#

```

Number of children	Frequency
0	0
1	1
2	2
3	3
4	10
5	8
6	5
7	3
8	2
9	1
10	1

[illegible]

Number of children	Frequency
0	0
1	1
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3	3
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6	5
7	3
8	2
9	1
10	1

[illegible]





Number of children	Frequency
0	0
1	1
2	2
3	1
4	10
5	2
6	1
7	2
8	1
9	3
10	4

[illegible]

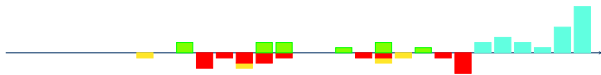




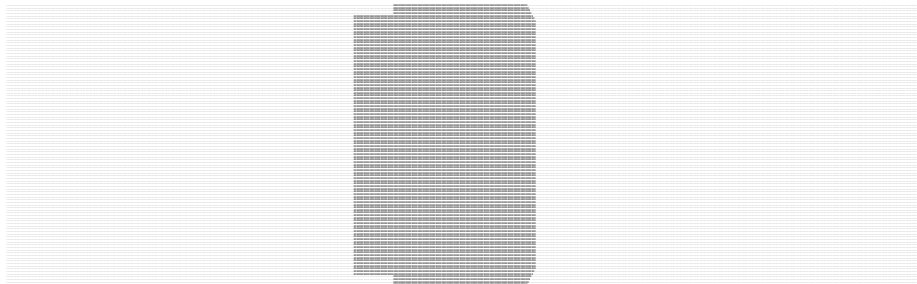


# G: Gluttonous Goop

Problem Author: Mees de Vries

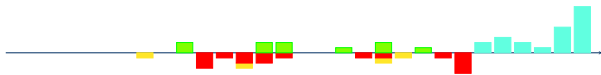


Eventually everything turns into: a big rectangle with weird corners.



# G: Gluttonous Goop

Problem Author: Mees de Vries



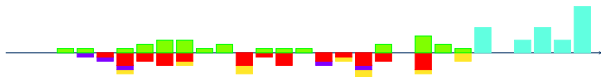
## Solution:

- Simulate for 20 steps, then compute the height/width of the bounding rectangle.
- Count how many squares in the rectangle are missing.
- Find the height  $\times$  width for the final bounding rectangle. Subtract corner squares.

Statistics: 52 submissions, 10 accepted, 22 unknown

# H: Historic Exhibition

Problem Author: Bruno Ploumhans



Given  $k$  vases, and  $p$  pedestals, place every vase on a pedestal of the right size.

Observations:

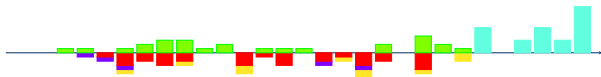
- If a pedestal has only one size, match those to vases of that size.
- Then you only have pedestals  $(1, 2), (2, 3), (3, 4), \dots$
- Match all 1-vases to  $(1, 2)$ -pedestals.
- Use leftover  $(1, 2)$ -pedestals for 2-vases.
- Then you only have pedestals  $(2, 3), (3, 4), (4, 5), \dots$
- ... and repeat.

In other words: go greedily from left to right.



# H: Historic Exhibition

Problem Author: Bruno Ploumhans



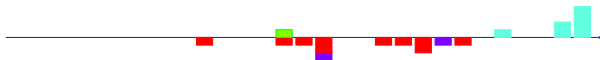
## Common pitfalls:

- Output impossible once the smallest vase does not fit the smallest available pedestal. (There might still be larger pedestals!)
- Flow algorithm: too slow.

Statistics: 100 submissions, 26 accepted, 29 unknown

# I: Inquiry II

Problem Author: Timon Knigge



- Given a graph, find a maximum independent set.
- Normally this is NP complete, but ... this graph is very close to a tree!
- We can solve this problem in linear time for a tree.
- Solution:
  - Find the  $k$  additional edges, at most 16.
  - For each additional edge, at least one end point is not in the independent set. Brute force all  $2^k$  options.
- Total runtime  $O(2^k \cdot n)$ .

Statistics: 20 submissions, 1 accepted, 7 unknown

# J: Jazz it Up!

Problem Author: Ragnar Groot Koerkamp



Given a squarefree number  $n$ , find an  $1 < m < n$  such that  $n \times m$  is squarefree.

- Squarefree test: trial division by all squares smaller than  $n$ .
- To find  $m$ : try all options starting at 2. This is fast since you'll always find a solution among the first 13 primes.
- Do **NOT** print  $n - 1$ . It fails for  $k^2 + 1$ .

Statistics: 80 submissions, 56 accepted, 0 unknown

# J: Jazz it Up!

Problem Author: Ragnar Groot Koerkamp



```
static int primes[]={2,3,5,7,11,13,17,19,23,29,31,37, ..., 99991}

int main() {
    int n;
    cin >> n;
    for (int i = 0; i < (int)sizeof(primes) / (int)sizeof(primes[0]);
        i++) {
        if (n % primes[i] != 0) {
            cout << primes[i] << endl;
            break;
        }
    }
}
```

# J: Jazz it Up!

Problem Author: Ragnar Groot Koerkamp



```
t = int(input())
P = [2,3,5,7,11,13,17,19,23,29]
for i in P:
    if t%i!=0:
        print(i)
        break
```

# J: Jazz it Up!

Problem Author: Ragnar Groot Koerkamp



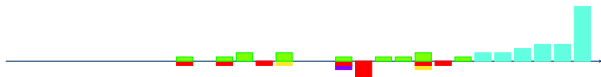
```
# Get input
n = int(raw_input())

# Find a number m such that m is relatively prime to n

# Check if any of these numbers divides n
if n%2 == 0:
    if n%3 == 0:
        if n%5 == 0:
            if n%7 == 0:
                if n%11 == 0:
                    if n%13 == 0:
                        if n%17 == 0:
                            print 19
                        else:
                            print 17
                    else:
                        print 13
                else:
                    print 11
            else:
                print 7
        else:
            print 5
    else:
        print 3
```

# K: Keep Him Inside

Problem Author: Timon Knigge



- The problem was to find a weighted average of the vertices of some convex polygon that equals the given point  $P$ .
- Choose one of the vertices of the polygon as a “base point”, translate so the base point is the origin and triangulate the polygon by drawing lines from the base point.
- Find the triangle in which the prisoner  $P$  lies (e.g. by calculating angles from the base point). Call the vectors from the base point to two other vertices of the triangle  $\mathbf{v}_1$  and  $\mathbf{v}_2$
- The vector  $P$  can now be decomposed into  $a \times \mathbf{v}_1 + b \times \mathbf{v}_2$  by projecting (calculate some inner products).
- The weights are  $1 - a - b$ ,  $a$  and  $b$  for the three points of the triangle and 0 otherwise.

Statistics: 53 submissions, 12 accepted, 28 unknown



## L: Lucky Draw

Problem Author: Raymond van Bommel and Mees de Vries

At the casino,  $n$  players start with  $k$  lives each. Each round, each one loses a life with probability  $1 - p$ . You win if you are the only one remaining. What is the probability of a draw?

$$\begin{aligned}\mathbb{P}(\text{draw}) &= 1 - \mathbb{P}(\text{someone wins}) \\ &= 1 - n \times \mathbb{P}(\text{player 1 wins}) \\ &= 1 - n \times \sum_{i=1}^{\infty} \mathbb{P}(\text{player 1 dies round } i, \text{ other player die before round } i) \\ &= 1 - n \times \sum_{i=1}^{\infty} \mathbb{P}(\text{player 1 dies round } i) \times \mathbb{P}(\text{Player 1 dies before round } i)^{n-1}.\end{aligned}$$

For large  $i$ ,  $\mathbb{P}(\text{player 1 dies round } i)$  is very small. So compute only for  $i$  up to  $M = 1000$  (or more).



# L: Lucky Draw

Problem Author: Raymond van Bommel and Mees de Vries

Two options:

1 Mathematically:

$$\mathbb{P}(\text{Player 1 dies round } i) = \binom{i-1}{k-1} p^{i-k} (1-p)^k.$$

This gives an  $\mathcal{O}(M)$  algorithm.

2 With dynamic programming: let

$$DP[r][l] = \mathbb{P}(\text{Player 1 has } l \text{ lives in round } r).$$

Then:

$$DP[r][l] = pDP[r-1][l] + (1-p)DP[r-1][l+1]$$

(plus correct edge conditions). This gives an  $\mathcal{O}(Mk)$  algorithm.

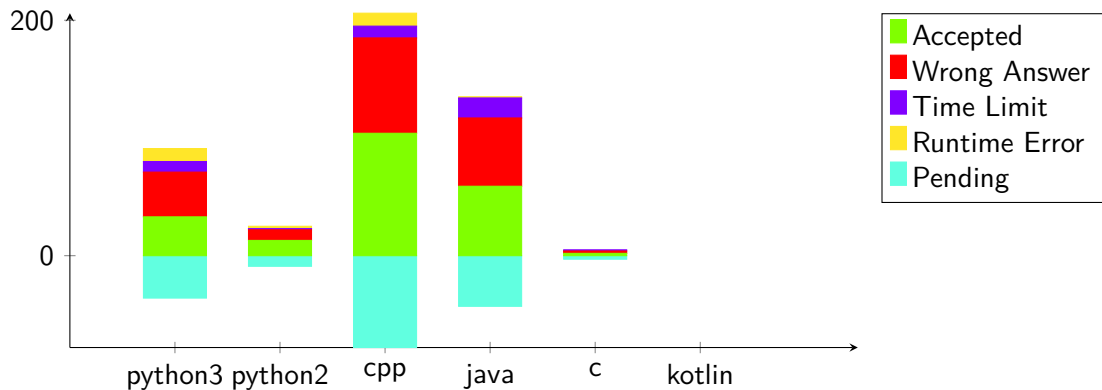
# L: Lucky Draw

Problem Author: Raymond van Bommel and Mees de Vries



Statistics: 3 submissions, 2 accepted, 0 unknown

## Language stats



# The Proofreaders

- Jelle Besseling
- Job Doesburg
- Nicky Gerritsen
- Raymond van Venetië
- Mees Vermeulen
- Jan Westerdiep

# The Jury

- Onno Berrevoets
- Daan van Gent
- Ragnar Groot Koerkamp
- Bjarki Ágúst Guðmundsson
- Joey Haas
- Timon Knigge
- Harry Smit
- David Venhoek
- Mees de Vries
- Wessel van Woerden