

NWERC 2010

Solutions to the problems

The Jury

Jacobs University Bremen

Problem A

Problem H

Problem C

Problem E

Problem B

Problem F

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Problem D

Problem I



A - Fair Division

- ▶ Sort persons according to maximum contribution
- ▶ Tie-breaker: position in list
- ▶ for (i=0 ... N-1)
- ▶ `contrib[i] = min(max[i] , price/(N-i))`
- ▶ `price -= contrib[i]`
- ▶ Don't print a trailing space

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A - Fair Division

- ▶ Sort persons according to maximum contribution
- ▶ Tie-breaker: position in list
- ▶ for (i=0 ... N-1)
- ▶ `contrib[i] = min(max[i] , price/(N-i))`
- ▶ `price -= contrib[i]`
- ▶ Don't print a trailing space

- ▶ Statistics: 119 submissions, 51 correct, first 27 minutes

Solutions

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H - Stock Prices

- ▶ While bid price larger than ask price, process deals
- ▶ Output prices or a dash if it doesn't exist

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H - Stock Prices

- ▶ While bid price larger than ask price, process deals
- ▶ Output prices or a dash if it doesn't exist

- ▶ Statistics: 112 submissions, 48 correct, first 40 minutes

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C - High Scores

- ▶ Loop over starting with *going left* or *right*
- ▶ Loop over where to *turn around*
- ▶ Count the number of moves until you are done

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C - High Scores

- ▶ Loop over starting with *going left* or *right*
- ▶ Loop over where to *turn around*
- ▶ Count the number of moves until you are done

- ▶ Statistics: 192 submissions, 43 correct, first 34 minutes

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E - Rankings

- ▶ Start with `newrank[i] = oldrank[i]`
- ▶ For a swap (i, j) , increase/decrease `newrank[i | j]`
- ▶ Check consistency: if i and j swapped, `newranks` and `oldranks` must be in opposite order
- ▶ There are never question marks in the answer
- ▶ Topological sorting also works

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E - Rankings

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- ▶ Check consistency: if i and j swapped, `newranks` and `oldranks` must be in opposite order
- ▶ There are never question marks in the answer
- ▶ Topological sorting also works

- ▶ Statistics: 77 submissions, 37 correct, first 65 minutes

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B - Free Goodies

- ▶ Sort the goodies to Petra's valuations
- ▶ $O(n^2)$ dynamic programming:
 - ▶ `best[n goodies taken][Jan took k]`
- ▶ Also $O(n \log n)$ greedy solution possible!

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B - Free Goodies

- ▶ Sort the goodies to Petra's valuations
- ▶ $O(n^2)$ dynamic programming:
 - ▶ `best[n goodies taken][Jan took k]`
- ▶ Also $O(n \log n)$ greedy solution possible!

- ▶ Statistics: 25 submissions, 9 correct, first 135 minutes

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F - Risk

- ▶ Binary search on the weakest link
- ▶ Use maximum flow algorithm to determine if answer is possible
- ▶ Graph vertices: source, sink, and 2 vertices for each land you control
- ▶ Graph edges:
 - ▶ source \rightarrow 1st land (cap=num. armies)
 - ▶ 1st land \rightarrow 2nd land (if connected)
 - ▶ 2nd land \rightarrow sink (cap=needed armies)

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F - Risk

- ▶ Binary search on the weakest link
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- ▶ Graph vertices: source, sink, and 2 vertices for each land you control
- ▶ Graph edges:
 - ▶ source \rightarrow 1st land (cap=num. armies)
 - ▶ 1st land \rightarrow 2nd land (if connected)
 - ▶ 2nd land \rightarrow sink (cap=needed armies)

- ▶ Statistics: 18 submissions, 8 correct, first 159 minutes

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J - Wormly

- ▶ Note: legs $2 \dots L - 1$ don't really matter
- ▶ Greedily move first leg, then last leg, then bubbles
- ▶ Repeat until finished
- ▶ Watch out for overflow

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J - Wormly

- ▶ Note: legs $2 \dots L - 1$ don't really matter
 - ▶ Greedily move first leg, then last leg, then bubbles
 - ▶ Repeat until finished
 - ▶ Watch out for overflow
-
- ▶ Statistics: 102 submissions, 9 correct, first 161 minutes

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G - Selling Land

- ▶ Process rows one by one
- ▶ For each column c of row r , count the number of grass squares above (c, r)
- ▶ Process columns and keep a list of possible end squares
- ▶ This takes amortized time $O(1)$ per square

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G - Selling Land

- ▶ Process rows one by one
 - ▶ For each column c of row r , count the number of grass squares above (c, r)
 - ▶ Process columns and keep a list of possible end squares
 - ▶ This takes amortized time $O(1)$ per square
-
- ▶ Statistics: 21 submissions, ≥ 1 correct, first 254 minutes

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D - Hill Driving

- ▶ Drive through the landscape with constant speed
- ▶ (*derive for two segments with equations, then use induction*)
- ▶ Binary search and check how much fuel is used
- ▶ Be careful:
 - ▶ Don't gain fuel when going downhill, but go faster
 - ▶ Don't exceed maximum speed
- ▶ Linear search possible too

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D - Hill Driving

- ▶ Drive through the landscape with constant speed
- ▶ (*derive for two segments with equations, then use induction*)
- ▶ Binary search and check how much fuel is used
- ▶ Be careful:
 - ▶ Don't gain fuel when going downhill, but go faster
 - ▶ Don't exceed maximum speed
- ▶ Linear search possible too

- ▶ Statistics: 28 submissions, ?? correct, first ?? minutes

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I - Telephone Network

- ▶ All sets of requests are possible, so we can add dummy requests to get bipartite graph with $\deg(v) = 1$ for all v
- ▶ Reduce this graph modulo 2^{n-1} to get a bipartite graph with $\deg(v) = 2$ for all v
- ▶ Split this graph in two graphs with all degrees 1 and you get two instances of the same problem with $n' = n/2$
- ▶ Solve recursively and construct solution

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I - Telephone Network

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 - ▶ Reduce this graph modulo 2^{n-1} to get a bipartite graph with $\deg(v) = 2$ for all v
 - ▶ Split this graph in two graphs with all degrees 1 and you get two instances of the same problem with $n' = n/2$
 - ▶ Solve recursively and construct solution
-
- ▶ Statistics: 6 submissions, ?? correct, first ?? minutes

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The end



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