



# Bergen Open 2019

## Solution Slides

November 2, 2019



UNIVERSITY OF BERGEN

# The Jury

- Olav Røthe Bakken
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Special thanks:



- Greg Hamerly (Kattis)
- Kirill Simonov (for testing problems)

# Howl



- Problem summary: Give a longer howl than Fenrir. Howl must follow given rules.
- Algorithms:
  - `print("A"*(len(input())) + "WHO")`
  - `print(input() + "O")`
  
- Runtime:  $O(n)$

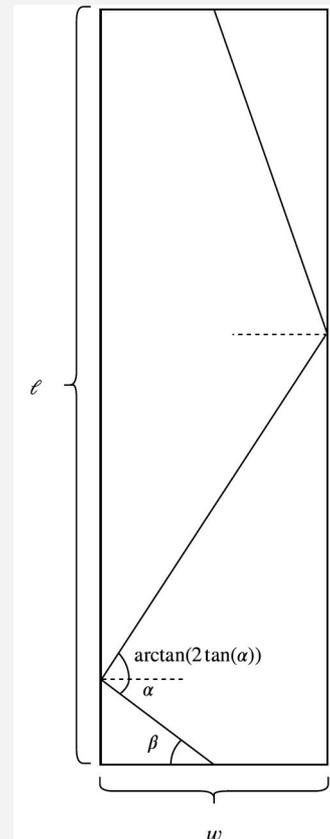
# Climbing stairs



- Problem summary: How many steps are we required to walk each day in order to participate in the staircase cup
- Algorithm:
  - We can always postpone registering to the last possible moment
  - Therefore we will first go to our office, then register at the end of day, then go home
  - If we don't have enough steps when we get to the registration office, pad the number of steps until we have enough, going two steps at a time
  - `print (max(n, k + abs(r-k)) + r + (1 if n%2 != r%2 and n > k + abs(k-r) else 0))`
- Runtime:  $O(1)$

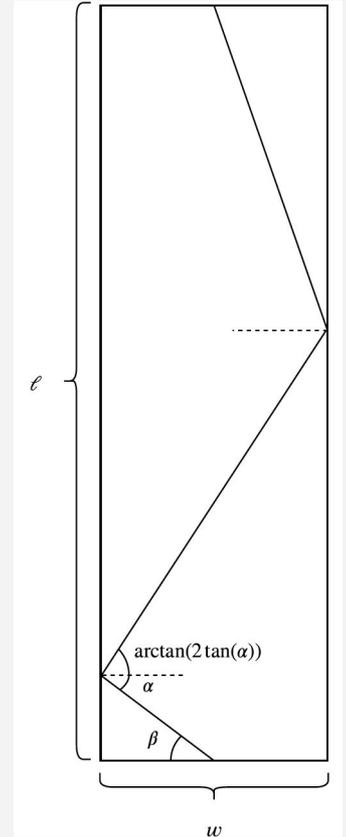
# Fence bowling

- Problem summary: Determine angle such that you hit strike after  $k$  bounces.
- Algorithm 1:
  - Binary search on angle  $\beta$
  - For each guess, simulate bounces
- Runtime:  $O(k \log(1/\epsilon))$



# Fence bowling

- Problem summary: Determine angle such that you hit strike after  $k$  bounces.
- Algorithm 2:
  - Observe: the triangle before a bounce is half as “long” (along the centre line) as the triangle after the bounce.
  - There are  $k$  pairs of right triangles (following the path from centre line, to side rail, back to centre line), each pair twice as long as the previous pair.
  - Let the first pair of triangles “stretch” a length  $x$ . Then total length  $L = x + 2x + \dots + 2^{k-1}x$
  - Hence,  $x = L / (2^k - 1)$
  - Answer is  $\arctan(L / (2^k - 1) / 3 / (w / 2))$
- Runtime:  $O(1)$



# Bus Ticket



- Problem summary: Decide when to buy single tickets and when to buy period tickets, such that the total cost is minimized.
- Dynamic programming
  - Create array  $dp[n]$
  - Define  $dp[i]$  to be minimum cost to purchase the trips  $0\dots i$
  - Base case:  $dp[0]$  is price of single ticket (or period ticket, if this is cheaper)
  - Recursive case:  $dp[i]$  is the minimum of
    - buying a single ticket for the last trip:  $dp[i-1] + \text{price of single trip}$
    - buying a period ticket for the last trip:  $dp[j] + \text{price of period ticket}$ , where  $j$  is the latest trip for which a period ticket can not cover both trip  $j$  and trip  $i$ .
- Runtime:  $O(n^2)$

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- Runtime:  ~~$\Theta(n^2)$~~   $O(n \log n)$  (with binary search to find  $j$ )

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- Runtime:  ~~$\Theta(n^2)$~~   ~~$\Theta(n \log n)$~~   $O(n)$  (with sliding pointer to find  $j$ )

**Author:** Torstein Strømme

**First solved:** 01:36

**Solved by:** 1 team

# Alehouse



- Problem summary: During which time interval of length  $k$  can you meet the most different people in the alehouse?
- What if interval has length 0?
  - For each person, make two events: Arrival and departure.
  - Sort all events (sort arrivals before departures)
  - count = 0
  - for each event in events:
    - if event is arrival, count++
    - if event is departure, count--
  - Remember maximum value of count

# Alehouse



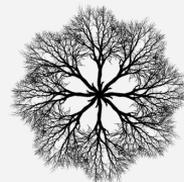
- Problem summary: During which time interval of length  $k$  can you meet the most different people in the alehouse?
- What if interval has length 0?
  - Can solve in time  $O(n \log n)$
- Observation:
  - You stay for  $k$  seconds  $\Leftrightarrow$  You stay for 0 seconds, everyone else stays for  $k$  seconds longer
  
- Runtime:  $O(n \log n)$

# Great GDP



- Problem summary: Find the connected subtree containing the root with the largest gdp per capita
- Algorithm:
  - If the root has the largest gdp per capita we can simply select the root
  - Otherwise some other vertex  $v$  has largest gdp per capita
  - Every solution which includes  $v$  must also include  $\text{parent}[v]$
  - Merge  $v$  and  $\text{parent}[v]$ 
    - Can use union-find to keep track of gdp, population and parent
  - Use a priority queue to quickly find the vertex with highest gdp per capita
- Runtime:  $O(n \log n)$

# Equilibrium



- Problem summary: Find the order of vertices which minimizes imbalance
- Algorithm:
  - There exists an ordering where every vertex with even degree has imbalance 0, and every vertex with odd degree has imbalance 1
  - Pick a vertex as the root, and distribute its children evenly on either side
  - Disjoint subtrees will not interfere with each other, so we can assume the vertices from each subtree are contiguous in the optimal ordering
  - Recursively find the order of each subtree
  
- Runtime:  $O(n)$

# Killing Chaos



- Problem summary: Figure out the maximum chaos according to the rules
- Rules: chaos = # of train segments \* sum(round up to closest 10 the # of passengers in each segment)
- Naive algorithm:
  - Simulate the process
  - Keep an array which keeps track of whether each wagon is killed
  - Each time a wagon is blown up, recalculate the chaos
- Runtime:  $O(n^2)$

# Killing Chaos



- Problem summary: Figure out the maximum chaos according to the rules
- Rules: chaos = # of train segments \* sum(round up to closest 10 of passengers in each segment)
- Better algorithm:
  - Simulate the process *backwards*
  - Use union-find to keep track of how many passengers in each segment
  - Keep track of number of segments, and the “base chaos” (before multiplication with number of segments)
- Runtime:  $O(n \log^* n)$

# Killing Chaos



- Problem summary: Figure out the maximum chaos according to the rules
- Rules: chaos = # of train segments \* sum(round up to closest 10 of passengers in each segment)
- Another good algorithm:
  - Keep a sorted set (binary search tree) which contains train segments (lower bound, upper bound, # of people)
  - Keep track of number of segments, and the “base chaos” (before multiplication with number of segments)
  - When a coach is killed, remove corresponding segment from sorted set (found in  $\log(n)$  time), and add back smaller segments if necessary.
- Runtime:  $O(n \log n)$

# Jane Eyre



- Problem summary: Given that Anna always reads in her books in alphabetical ASCII order, when will she (at the earliest) finish reading Jane Eyre? Books arrive as time goes.
- Simulation
  - Let time be 0
  - Pick the earliest book from priority queue sorted by ASCII order; read it and update time
  - Receive all new books that arrive at current time or earlier, put those in priority queue (use sliding pointer)
  - Repeat until Jane Eyre is read
- Runtime:  $O(n \log n)$

# Jane Eyre



- Problem summary: Given that Anna always reads in her books in alphabetical ASCII order, when will she (at the earliest) finish reading Jane Eyre? Books arrive as time goes.
- Alternative simulation
  - Ignore all books after Jane Eyre in ASCII alphabet
  - Sort books by arrival time
  - Read the books, track the time; continue until the next book arrives after the current time
  - Return current time + time to read Jane Eyre
- Runtime:  $O(n \log n)$

# Ice cream



- Problem summary: Produce as much chocolate ice cream as possible
- Algorithm:
  - We want to compute the maximum amount of flow ( $W$ ) from  $c$  and  $v$  to  $f$ , such that the flow from the chocolate tank  $c$  is equal to the flow from the vanilla tank  $v$ .
  - Convert into a standard max flow problem by binary search for the answer
    - Add a super-source with pipes to  $c$  and  $v$  that each have capacity  $g$  (half the guessed flow)
    - It is possible the optimal solution uses half integral amounts of each ingredient
  - Implement using your favourite max-flow algorithm (e. g. Edmund's Karp)
- Runtime:  $O(nm^2 \log W)$

# Drive safely



- Problem summary: Given a polyline describing a road, place speed signs such that travel time by travelling legally is as small as possible.
- Some basic geometry to find angles and distances
- Dynamic programming:
  - Two tables:  $dp\_a[n][k]$  and  $dp\_b[n][k]$
  - Define  $dp\_a[i][j]$  = Minimum time required to travel to (just before) point  $i$  using  $j$  or less speed signs
  - Define  $dp\_b[i][j]$  = Minimum time required to travel to (just after) point  $i$  using  $j$  or less speed signs
  - At location  $i$ , check every possible location for the previous speed sign
- Runtime:  $O(n^2k)$

# Statistics

- Number of teams: 12
- Number of participants: 30
- Number of submissions: 180
- Number of accepted submissions: 35
- First accepted submission: 00:07:54 (Howl)
- Last accepted submission: 04:51:02 (Jane Eyre)
- Number of commits to problem repository: 164

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